



Victoria University
of Wellington, New Zealand
*Te Whare Wananga o te
Upoko o te Ika a Maui
Aotearoa*

The Whiley Programming Language

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<http://whiley.org>

Java:

```
void buildLabelMap(List<Bytecode> bytecodes) {
    HashMap<String,Integer> labels = new HashMap<String,Integer>();
    int idx = 0
    for(Bytecode b : bytecodes) {
        if(b instanceof Bytecode.Label) {
            Bytecode.Label lab = (Bytecode.Label) b;
            labels.put(lab.name, idx);
        }
        idx = idx + 1;
    } }
```

Python:

```
def buildLabelMap(bytecodes):
    labels = {}
    idx = 0
    for b in bytecodes:
        if type(b) == "Label":
            labels[b.name] = idx
    idx = idx + 1
```

The Whiley Language

```
void buildLabelMap([Bytecode] bytecodes):  
  labels = {->}  
  idx = 0  
  for b in bytecodes:  
    if b is Label:  
      labels[b.name] = idx  
  idx = idx + 1
```

- Design Goals:
 - **Look and feel** of a dynamically typed language
 - But, still provide **static** type checking
 - **Simple** programming model
 - Amenable to **program verification**

Data Types

- **Unbound Integers and Rationals**
 - `int` currently implemented `BigInteger`
 - `real` currently implemented as `BigRational`
 - How to do this efficiently on JVM?
- **Lists, Sets and Maps:**

```
int sum([int] list):  
    r = 0  
    for x in list:  
        r = r + x  
    return r
```

- **Records:**

```
define Point as {int x, int y}
```

Constraints

```
int sum([int] list) requires no {x in $ | x < 0},  
                    ensures $ >= 0:  
  
  r = 0  
  for x in list:  
    r = r + x  
  return r
```

```
define nat as int where $ >= 0  
define natlist as [nat]  
  
nat sum(natlist list):  
  ...  
  return r
```

- Pre/Post Conditions and Invariants:
 - Goal is to check at compile time
 - Currently, checked at run time

Flow-Sensitive Typing

```
Value evaluate(Expr e, {string->Value} env) throws Error:  
  if e is int:  
    return e  
  else if e is Var:  
    return env[e.id]  
  else if e is BinOp:  
    lhs = evaluate(e.lhs, env)  
    rhs = evaluate(e.rhs, env)  
  ...
```

- Flow-Sensitive Typing:
 - Variables declared by **assignment**
 - Variables automatically **retyped** by type tests
 - Variables can have **different types** at different points

Union Types

```
nullable|int indexOf(string str, char c):
    ...

[string] split(string str, char c):
    idx = indexOf(str,c)
    if idx is int:
        below = str[0..idx]
        above = str[idx..]
        return [below,above]
    else:
        return [str]
```

- Union types have many uses
 - Such as neatly handling null references
 - Or, combining different kinds (i.e. unions of structs)

Structural Subtyping

```
define LinkedList as null | {int dat, LinkedList nxt}
```

```
int sum(LinkedList l):
```

```
  if l == null:
```

```
    return 0
```

```
  else:
```

```
    return l.dat + sum(l.nxt)
```

```
void main(System sys, [string] args):
```

```
  l={dat: 1, nxt: null}
```

```
  l={dat: 2, nxt: l}
```

```
  sys.out.println(sum(l))
```

- Defined types are **not nominal**
 - i.e. LinkedList is just a name that "expands"

Structural Subtyping

```
define LinkedList as null | {int dat, LinkedList nxt}
```

```
int sum(LinkedList l):
```

```
  if l == null:
```

```
    return 0
```

```
  else:
```

```
    return l.dat + sum(l.nxt)
```

```
void main(System sys, [string] args):
```

```
  l={dat: 1,nxt: null} // l is {int dat, null nxt}
```

```
  l={dat: 2,nxt: l} // l is {int dat, {int dat, null nxt} nxt}}
```

```
  sys.out.println(sum(l))
```

- Defined types are **not nominal**
 - i.e. LinkedList is just a name that "expands"

Value Semantics

```
Board applyMove(Move move, Board board) throws Move.Invalid:  
    nboard = applyMoveDispatch(move,board)  
    if !validMove(move,board,nboard):  
        throw Move.Invalid(board,move)  
    else:  
        return nboard
```

- **Whiley does not have references!**
 - **Everything** is pass-by-value
 - Data propagates only via return
 - Much more **functional** in nature
 - Requires different way of **thinking**

Reference Counting

```
Board applySimple(Board b, Pos old, Pos pos, Piece piece):  
    b[old.col][old.row] = null  
    b[pos.col][pos.row] = piece  
    return board
```

```
Board applyMove(Board b, Move m):  
    if move is Simple:  
        ...  
        return update(b, m.from, m.to, m.piece)
```

- Value semantics:
 - Copy board for call to applySimple()
 - Copy again for assignments in applySimple()
 - But, this is very inefficient!!!
 - Reference counting can really help here...

Digression

- Item 24, *Effective Java*
 - Make Defensive Copies when Needed

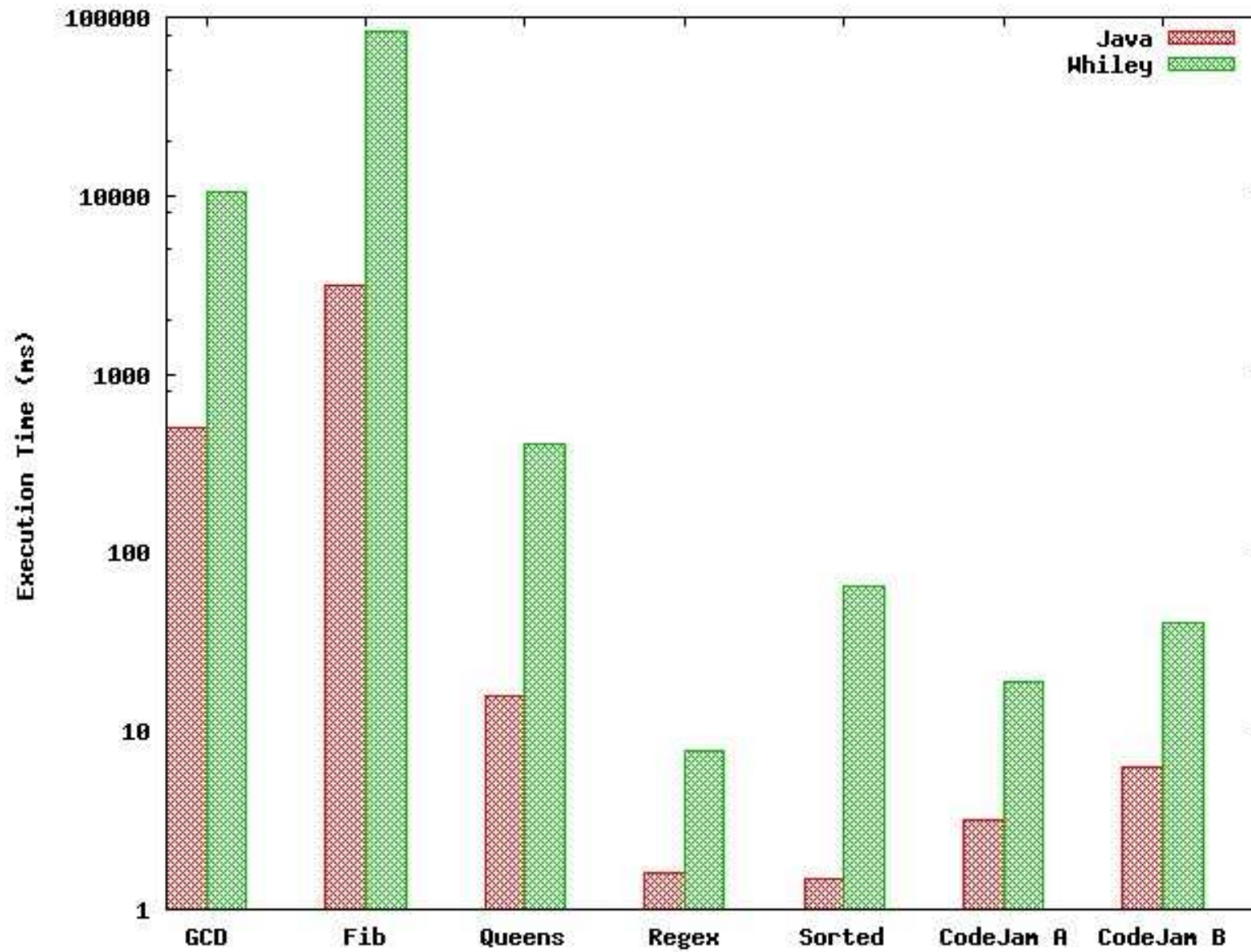
“It is essential to make a defensive copy of each mutable parameter to the constructor”

-- Josh Bloch

Effectiveness of Ref Counting

Benchmark	LOC	# Clones	% Clones
Gunzip	815	873 / 140561	0.62%
JASM	2333	12878 / 29968	43.0%
Chess	784	6438 / 416116	1.6%
Calc	225	0 / 81527	0.0%
SCC	169	12602 / 258968	4.86%

Performance



Java Interoperation

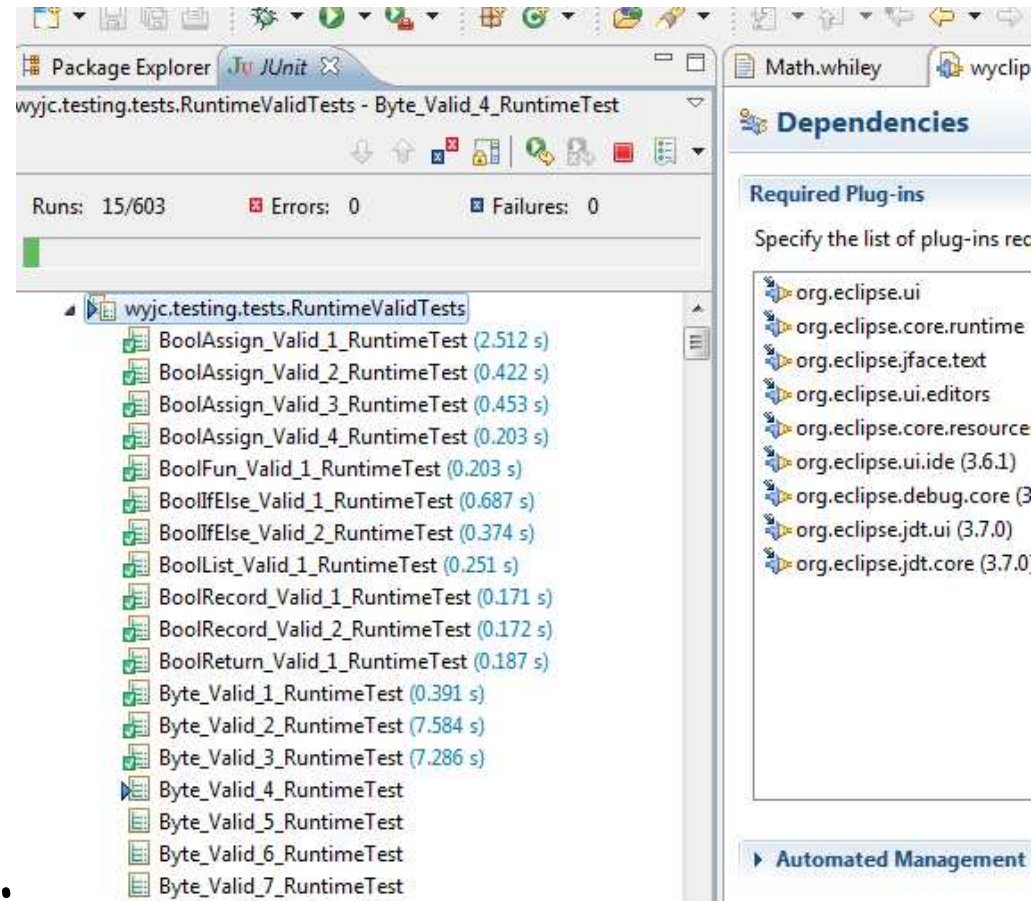
Sum.wiley:

```
...  
  
public native int sum([int] list):  
  
...
```

Sum\$native.java:

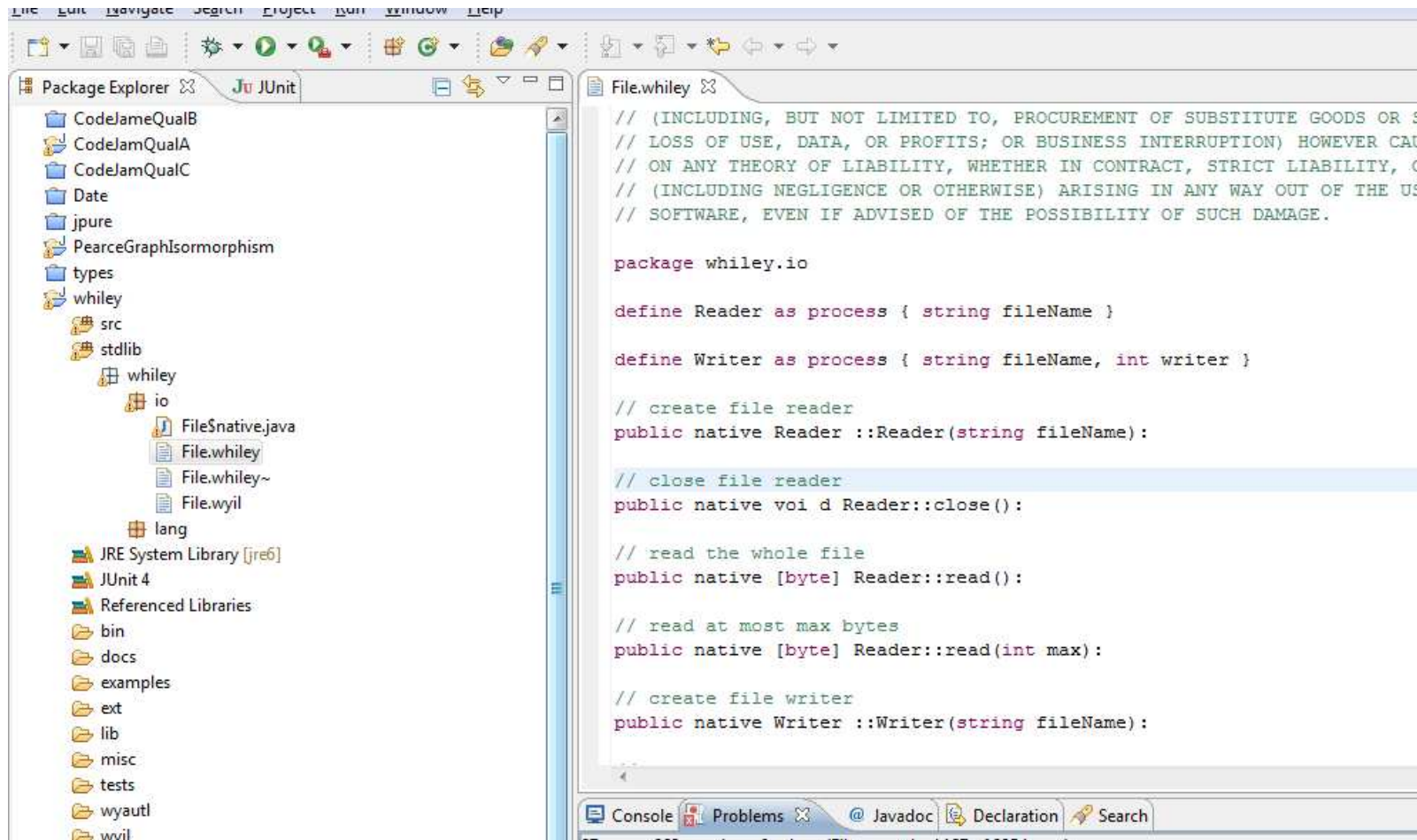
```
...  
  
public BigInteger sum(wyjc.runtime.List list) {  
    ...  
}  
  
...
```

Testing



- Testing:
 - Some 600 end-end tests
 - Over 15,000 unit tests of type system!

Eclipse Plugin



- Update Site: <http://whiley.org/eclipse>



<http://whiley.org>